Theorem The Erlang distribution has the scaling property. That is, if $X \sim \text{Erlang}(\alpha, n)$ then Y = kX also has the Erlang distribution.

Proof Let the random variable X have the $Erlang(\alpha, n)$ distribution with probability density function

$$f(x) = \frac{x^{n-1}e^{-x/\alpha}}{\alpha^n(n-1)!} \qquad x > 0.$$

Let k be a positive, real constant. The transformation Y = g(X) = kX is a 1–1 transformation from $\mathcal{X} = \{x \mid x > 0\}$ to $\mathcal{Y} = \{y \mid y > 0\}$ with inverse $X = g^{-1}(Y) = Y/k$ and Jacobian

$$\frac{dX}{dY} = \frac{1}{k}.$$

Therefore, by the transformation technique, the probability density function of Y is

$$f_Y(y) = f_X(g^{-1}(y)) \left| \frac{dx}{dy} \right|$$

$$= \frac{(y/k)^{n-1} e^{-y/k\alpha}}{\alpha^n (n-1)!} \left| \frac{1}{k} \right|$$

$$= \frac{y^{n-1} e^{-y/k\alpha}}{(k\alpha)^n (n-1)!} \qquad y > 0,$$

which is the probability density function of a Erlang $(k\alpha, n)$ random variable.

APPL verification: The APPL statements

```
assume(k > 0);
X := [[x -> x ^ (n - 1) * exp(-x / alpha) / (alpha ^ n * factorial(n - 1))],
        [0, infinity], ["Continuous", "PDF"]];
g := [[x -> k * x], [0, infinity]];
Y := Transform(X, g);
```

yield the probability density function of a $Erlang(k\alpha, n)$ random variable.